## Homework 6

1. RSA Assumption $(\mathbf{5}+\mathbf{1 2 + 5})$. Consider RSA encryption scheme with parameters $N=35=5 \times 7$.
(a) Find $\varphi(N)$ and $\mathbb{Z}_{N}^{*}$.
(b) Use repeated squaring and complete the rows $X, X^{2}, X^{4}$ for all $X \in \mathbb{Z}_{N}^{*}$ as you have seen in the class (slides), that is, fill in the following table by adding as many columns as needed.
Solution.

| $X$ |  |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $X^{2}$ |  |  |  |  |  |  |  |  |  |  |  |  |
| $X^{4}$ |  |  |  |  |  |  |  |  |  |  |  |  |


| $X$ |  |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $X^{2}$ |  |  |  |  |  |  |  |  |  |  |  |  |
| $X^{4}$ |  |  |  |  |  |  |  |  |  |  |  |  |

(c) Find the row $X^{5}$ and show that $X^{5}$ is a bijection from $\mathbb{Z}_{N}^{*}$ to $\mathbb{Z}_{N}^{*}$. Solution.

| $X$ |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $X^{4}$ |  |  |  |  |  |  |  |  |  |  |


| $X$ |  |  |  |  |  |  |  |  |  |  |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $X^{4}$ |  |  |  |  |  |  |  |  |  |  |

2. Answer to the following questions $(7+7+7+7)$ :
(a) Compute the three least significant (decimal) digits of $87341011^{324562002}$ by hand. Solution.
(b) Is the following RSA signature scheme valid?(Justify your answer)
$(r \| m)=342454323, \sigma=13245345356, N=155, e=664$
Here, $m$ denotes the message, and $r$ denotes the randomness used to sign $m$ and $\sigma$ denotes the signature. Moreover, $(r \| m)$ denotes the concatenation of $r$ and $m$. The signature algorithm $\operatorname{Sign}(m)$ returns $(r \| m)^{d} \bmod N$ where $d$ is the inverse of $e$ modulo $\varphi(N)$. The verification algorithm $\operatorname{Ver}(m, \sigma)$ returns $\left((r \| m)==\sigma^{e}\right.$ $\bmod N)$.
Solution.
(c) Remember that in RSA encryption and signature schemes, $N=p \times q$ where $p$ and $q$ are two large primes. Show that in a RSA scheme (with public parameters $N$ and $e$ ), if you know $N$ and $\varphi(N)$, then you can find the factorization of $N$ i.e. you can find $p$ and $q$.

## Solution.

(d) Consider an encryption scheme where $\operatorname{Enc}(m):=m^{e} \bmod N$ where $e$ is a positive integer relatively prime to $\varphi(N)$ and $\operatorname{Dec}(c):=c^{d} \bmod N$ where $d$ is the inverse of $e$ modulo $\varphi(N)$. Show that in this encryption scheme, if you know the encryption of $m_{1}$ and the encryption of $m_{2}$, then you can find the encryption of $\left(m_{1} \times m_{2}\right)^{5}$.
Solution.

## Collaborators :

